

# Phonological Phrasing in Japanese

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SPOT at LSA 2021

\*Based on Bellik, Ito, Kalivoda, & Mester (to appear 2021)

# Japanese Mismatch

Kubozono (1989) found that a left-branching 4-word XP in Japanese maps to mismatching prosody:

- (1) [[[Naomi-no] ane-to] yunomi-no] iro]  
Naomi-GEN sister-GEN teacup-GEN color  
'the color of the teacup of Naomi's sister'

→ (<sub>φ</sub> (<sub>φ</sub> Naomi-no ane-no) (<sub>φ</sub> yunomi-no iro))

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*not a syntactic constituent*

*Evidence:* φ-initial rise on  $\omega_1, \omega_3$

Second rise is downstepped due to  $\phi^{\text{Max}}$  over (<sub>φ</sub>  $\omega_1 \omega_2$ ) and (<sub>φ</sub>  $\omega_3 \omega_4$ )

# Japanese Matches

However, 4-word XPs of all other shapes undergo perfect matching (Kubozono 1989):

*Right-branching:*

[A [B [C D]]] → (A (B (C D)))

*Balanced:*

[[A B] [C D]] → ((A B) (C D))

*Mixed (Left/Right):*

[[A [B C]] D] → ((A (B C)) D)

*Mixed (Right/Left):*

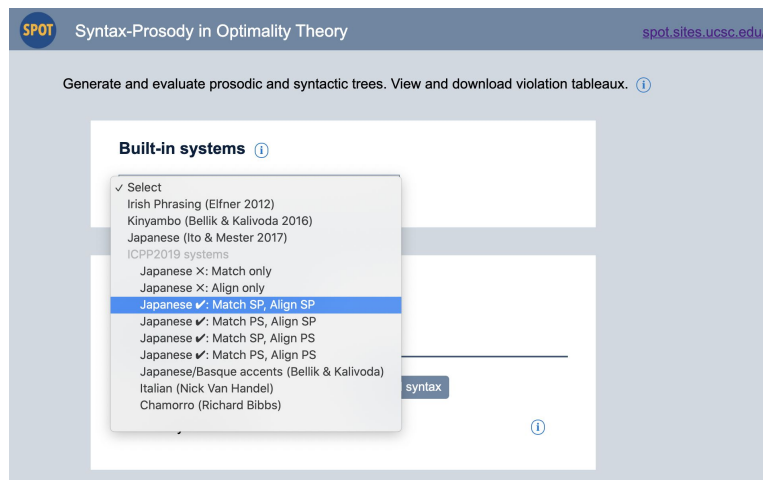
[A [[B C] D]] → (A ((B C) D))

Previous analyses (Selkirk 2011, Ishihara 2014, Kalivoda 2018) have attempted to analyze the left-branching mismatch in Match Theory (Selkirk 2011), but have not considered the matching cases.

We show that we need **Match and Align** to account for all these cases.

# Studying *OT systems*

- An OT system  $S = (\text{Gen}_S, \text{Con}_S)$
- We define OT systems by using **SPOT** (Bellik et al. 2015-2020) and **OTWorkplace** (Prince et al. 2007-2020).
- The systems discussed in this talk are on the SPOT interface (linked from <http://spot.sites.ucsc.edu>):



# Naming schema for our systems

S      'system'

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Msp     Match( $XP, \varphi$ ) in CON

Mps     Match( $\varphi, XP$ ) in CON



# Naming schema for our systems

S        'system'

Msp      Match(XP, $\varphi$ ) in CON

Mps      Match( $\varphi$ ,XP) in CON

Asp      Align(XP,L, $\varphi$ ,L) and Align(XP,R, $\varphi$ ,R) in CON

Aps      Align( $\varphi$ ,L,XP,L) and Align( $\varphi$ ,R,XP,R) in CON

# Naming schema for our systems

S        'system'

Msp      Match(XP, $\varphi$ ) in CON

Mps      Match( $\varphi$ ,XP) in CON

Asp      Align(XP,L, $\varphi$ ,L) and Align(XP,R, $\varphi$ ,R) in CON

Aps      Align( $\varphi$ ,L,XP,L) and Align( $\varphi$ ,R,XP,R) in CON

Only CON varies; GEN constant across systems.

# **S.Msp.Asp**

Matching and Alignment

# GEN.Msp.Asp: Inputs

A candidate is an input-output pair.

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## (1) *Inputs*

Syntactic trees with 3 or 4 terminal nodes, where:

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- every terminal node is an  $X^0$

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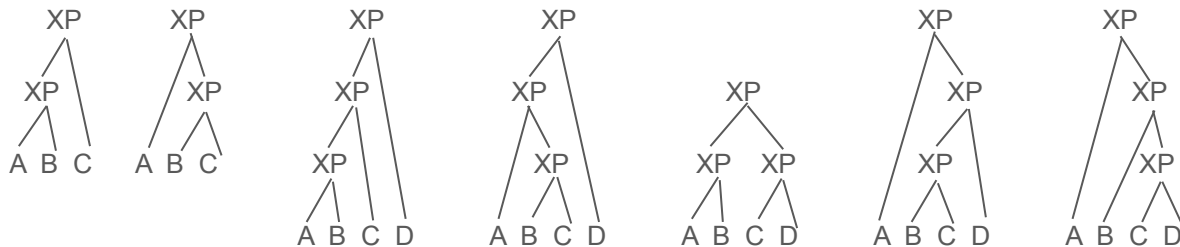
A candidate is an input-output pair.

## (1) *Inputs*

Syntactic trees with 3 or 4 terminal nodes, where:

- every non-terminal node is a binary-branching XP
- every terminal node is an  $X^0$

I.e.:





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For a syntactic input sTree, every prosodic tree pTree such that:

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For a syntactic input sTree, every prosodic tree pTree such that:

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- terminal nodes are of category  $\omega$

# GEN.Msp.Asp: Outputs

## (1) *Outputs*

For a syntactic input sTree, every prosodic tree pTree such that:

- non-terminal nodes are of category  $\varphi$
- terminal nodes are of category  $\omega$
- the terminal nodes in sTree stand in a one-to-one correspondence relation with the terminal nodes in pTree, with linear order preserved.

# CON.Msp.Asp

## (1) *Mapping constraints*

### (a) **MATCH(XP, $\varphi$ )**

Assign one violation for every node of category XP in the syntactic tree such that there is no node of category  $\varphi$  in the prosodic tree that dominates all and only the same terminal nodes as XP.

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Assign one violation for every node of category XP in the syntactic tree whose left edge is not aligned with the left edge of a node of category  $\varphi$  in the prosodic tree.

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### (c) **ALIGNR(XP, $\varphi$ )**

Assign one violation for every node of category XP in the syntactic tree whose right edge is not aligned with the right edge of a node of category  $\varphi$  in the prosodic tree.

# CON.Msp.Asp

## (1) *Mapping constraints*

- (a) MATCH( $XP, \varphi$ )
- (b) ALIGNL( $XP, \varphi$ )
- (c) ALIGNR( $XP, \varphi$ )

## (2) *Markedness constraints*

### (a) **BinMin( $\varphi, \omega$ )**

Assign one violation for every node of category  $\varphi$  in the prosodic tree that contains fewer than two nodes of category  $\omega$ .



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- (a) MATCH( $XP, \varphi$ )
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### (b) **BinMax( $\varphi, \omega$ )**

Assign one violation for every node of category  $\varphi$  in the prosodic tree that dominates more than two nodes of category  $\omega$ .

# CON.Msp.Asp

## (1) *Mapping constraints*

- (a) MATCH( $\text{XP}, \varphi$ )
- (b) ALIGNL( $\text{XP}, \varphi$ )
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## (2) *Markedness constraints*

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### (b) **BinMax( $\varphi, \omega$ )**

Assign one violation for every node of category  $\varphi$  in the prosodic tree that dominates more than two nodes of category  $\omega$ .

### (c) **BinMax( $\varphi, \text{branches}$ )**

Assign one violation for every node of category  $\varphi$  in the prosodic tree that has more than two children.

# Factorial Typology of S.Msp.Asp

	$[[[A \bar{B}] C] \bar{D}]$	$[[A [\bar{B} C]] D]$	$[A [[B C] D]]$	$[A [B [C D]]]$
L.1	$((A B) (C D))$	$((A B) (C D))$	$((A B) (C D))$	$((A B) (C D))$
L.2	$((A B) ((C) D))$	$((A B) ((C) D))$	$((A B) ((C) D))$	$((A B) (C D))$
L.3	$((A B) C) D$	$((A (B C)) D)$	$(A ((B C) D))$	$((A B) (C D))$
L.4	$((A B) C) D$	$(A (B C) D)$	$(A (B C) D)$	$((A B) (C D))$
L.5	$((A B) (C D))$	$(A (B C) D)$	$(A (B C) D)$	$((A B) (C D))$
L.6	$((A B) ((C) D))$	$(A (B C) D)$	$(A (B C) D)$	$((A B) (C D))$
L.7	$((A B) (C D))$	$((A (B)) (C D))$	$((A (B)) (C D))$	$((A (B)) (C D))$
L.8	$((A B) ((C) D))$	$((A (B)) ((C) D))$	$((A (B)) ((C) D))$	$((A (B)) (C D))$
L.9	$((A B) (C D))$	$(A (B C) D)$	$(A (B C) D)$	$((A (B)) (C D))$
L.10	$((A B) ((C) D))$	$(A (B C) D)$	$(A (B C) D)$	$((A (B)) (C D))$
L.11	$((A B) C) D$	$((A (B C)) D)$	$(A ((B C) D))$	$(A (B (C D)))$
<b>L.12</b>	<b><math>((A B) (C D))</math></b>	<b><math>((A (B C)) D)</math></b>	<b><math>(A ((B C) D))</math></b>	<b><math>(A (B (C D)))</math></b>
L.13	$((A B) C) D$	$(A (B C) D)$	$(A (B C) D)$	$(A (B (C D)))$
L.14	$((A B) (C D))$	$(A (B C) D)$	$(A (B C) D)$	$(A (B (C D)))$

# Factorial Typology of S.Msp.Asp

	[[[A B] C] D]	[[A [B C]] D]	[A [[B C] D]]	[A [B [C D]]]
L.1	((A B) (C D))	((A B) (C D))	((A B) (C D))	((A B) (C D))
L.2	((A B) ((C) D))	((A B) ((C) D))	((A B) ((C) D))	((A B) (C D))
L.3	(((A B) C) D)	((A (B C)) D)	(A ((B C) D))	((A B) (C D))
L.4	(((A B) C) D)	(A (B C) D)	(A (B C) D)	((A B) (C D))
L.5	((A B) (C D))	(A (B C) D)	(A (B C) D)	((A B) (C D))
L.6	((A B) ((C) D))	(A (B C) D)	(A (B C) D)	((A B) (C D))
L.7	((A B) (C D))	((A (B)) (C D))	((A (B)) (C D))	((A (B)) (C D))
L.8	((A B) ((C) D))	((A (B)) ((C) D))	((A (B)) ((C) D))	((A (B)) (C D))
L.9	((A B) (C D))	(A (B C) D)	(A (B C) D)	((A (B)) (C D))
L.10	((A B) ((C) D))	(A (B C) D)	(A (B C) D)	((A (B)) (C D))
L.11	(((A B) C) D)	((A (B C)) D)	(A ((B C) D))	(A (B (C D)))
<b>L.12</b>	<b>((A B) (C D))</b>	<b>((A (B C)) D)</b>	<b>(A ((B C) D))</b>	<b>(A (B (C D)))</b>
L.13	(((A B) C) D)	(A (B C) D)	(A (B C) D)	(A (B (C D)))
L.14	((A B) (C D))	(A (B C) D)	(A (B C) D)	(A (B (C D)))

*Not shown:*

[[A B] C]  
[A [B C]]  
[[A B] [C D]]

*These match in  
all 14 languages*

# Japanese pattern in S.Msp.Asp

	$[[[A \bar{B}] C] \bar{D}]$	$[[A [\bar{B} C]] D]$	$[A [[B C] D]]$	$[A [B [C D]]]$
L.1	$((A B) (C D))$	$((A B) (C D))$	$((A B) (C D))$	$((A B) (C D))$
L.2	$((A B) ((C) D))$	$((A B) ((C) D))$	$((A B) ((C) D))$	$((A B) (C D))$
L.3	$((A B) C) D$	$((A (B C)) D)$	$(A ((B C) D))$	$((A B) (C D))$
L.4	$((A B) C) D$	$(A (B C) D)$	$(A (B C) D)$	$((A B) (C D))$
L.5	$((A B) (C D))$	$(A (B C) D)$	$(A (B C) D)$	$((A B) (C D))$
L.6	$((A B) ((C) D))$	$(A (B C) D)$	$(A (B C) D)$	$((A B) (C D))$
L.7	$((A B) (C D))$	$((A (B)) (C D))$	$((A (B)) (C D))$	$((A (B)) (C D))$
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L.13	$((A B) C) D$	$(A (B C) D)$	$(A (B C) D)$	$(A (B (C D)))$
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← Japanese

# Japanese pattern in S.Msp.Asp

	$[[[A \bar{B}] C] \bar{D}]$	$[[A [\bar{B} \bar{C}]] D]$	$[A [[B C] D]]$	$[A [B [C D]]]$
L.1	$((A B) (C D))$	$((A B) (C D))$	$((A B) (C D))$	$((A B) (C D))$
L.2	$((A B) ((C) D))$	$((A B) ((C) D))$	$((A B) ((C) D))$	$((A B) (C D))$
L.3	$((A B) C) D$	$((A (B C)) D)$	$(A ((B C) D))$	$((A B) (C D))$
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L.5	$((A B) (C D))$	$(A (B C) D)$	$(A (B C) D)$	$((A B) (C D))$
L.6	$((A B) ((C) D))$	$(A (B C) D)$	$(A (B C) D)$	$((A B) (C D))$
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L.13	$((A B) C) D$	$(A (B C) D)$	$(A (B C) D)$	$(A (B (C D)))$
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Isomorphic mappings

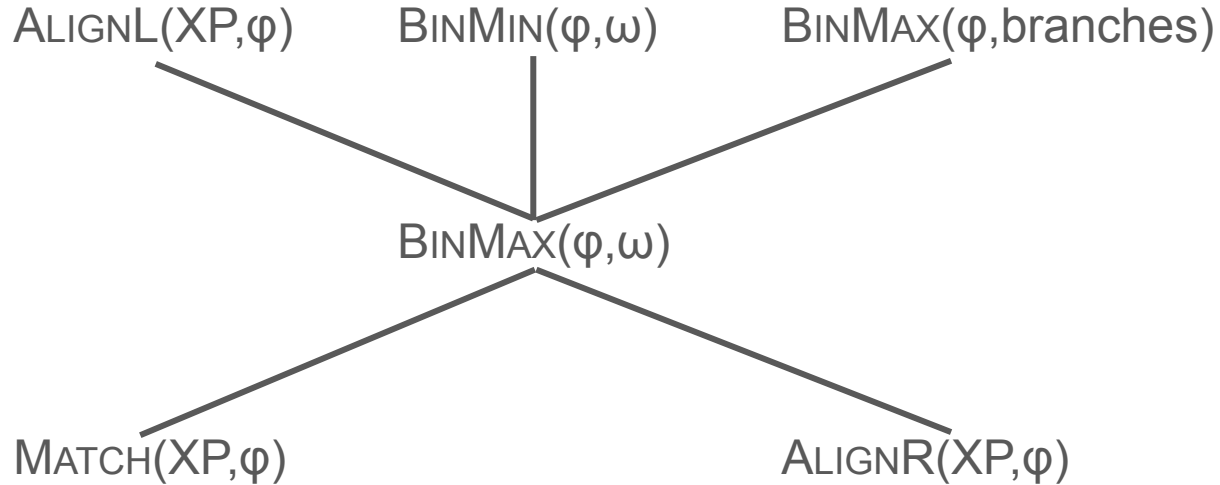
← Japanese

# Japanese pattern in S.Msp.Asp

	$[[[A \bar{B}] C] \bar{D}]$	$[[A [\bar{B} C]] D]$	$[A [[B C] D]]$	$[A [B [C D]]]$
L.1	$((A B) (C D))$	$((A B) (C D))$	$((A B) (C D))$	$((A B) (C D))$
L.2	$((A B) ((C) D))$	$((A B) ((C) D))$	$((A B) ((C) D))$	$((A B) (C D))$
L.3	$((A B) C) D$	$((A (B C)) D)$	$(A ((B C) D))$	$((A B) (C D))$
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L.8	$((A B) ((C) D))$	$((A (B)) ((C) D))$	$((A (B)) ((C) D))$	$((A (B)) (C D))$
L.9	$((A B) (C D))$	$(A (B C) D)$	$(A (B C) D)$	$((A (B)) (C D))$
L.10	<b>Rebracketing</b>	$(A (B C) D)$	$(A (B C) D)$	$((A (B)) (C D))$
L.11	$((A B) C) D$	$((A (B C)) D)$	$(A ((B C) D))$	$(A (B (C D)))$
<b>L.12</b>	<b><math>((A B) (C D))</math></b>	<b><math>((A (B C)) D)</math></b>	<b><math>(A ((B C) D))</math></b>	<b><math>(A (B (C D)))</math></b>
L.13	$((A B) C) D$	$(A (B C) D)$	$(A (B C) D)$	$(A (B (C D)))$
L.14	$((A B) (C D))$	$(A (B C) D)$	$(A (B C) D)$	$(A (B (C D)))$

← Japanese

# Grammar of L.12 in Sp.Msp.Asp

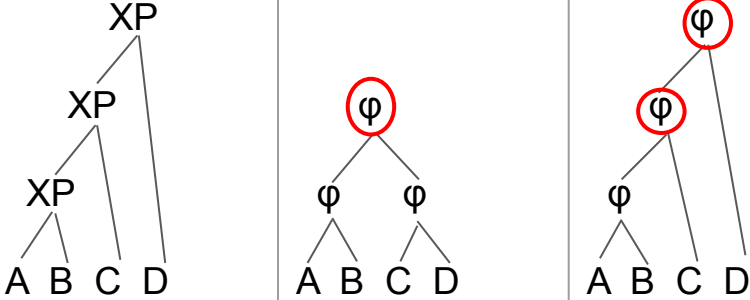




# Support for L.12

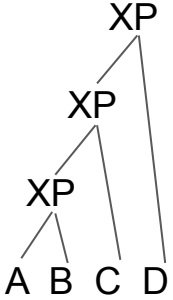
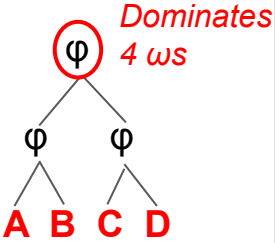
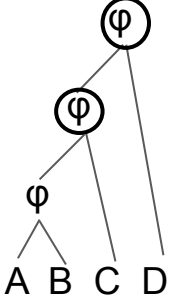
Input	Winner	Loser	ALIGNL (XP, $\varphi$ )	BINMIN ( $\varphi,\omega$ )	BINMAX ( $\varphi,br$ )	BINMAX ( $\varphi,\omega$ )	MATCH (XP, $\varphi$ )	ALIGNR (XP, $\varphi$ )
[A[B[CD]]]	(A(B(CD)))	((AB)(CD))	<b>W</b>			<b>L</b>	W	
[A[B[CD]]]	(A(B(CD)))	((A(B))(CD))		<b>W</b>		<b>L</b>	W	
[A[[BC]D]]	(A((BC)D))	(A(BC)D)			<b>W</b>	<b>L</b>	W	
[[[AB]C]D]	((AB)(CD))	((((AB)C)D))				<b>W</b>	<b>L</b>	<b>L</b>

# Support for L.12: *L-branching* $\rightarrow$ *rebracketed*

Input	Winner	Loser	ALIGNL (XP, $\varphi$ )	BINMIN ( $\varphi,\omega$ )	BINMAX ( $\varphi,br$ )	<b>BINMAX (<math>\varphi,\omega</math>)</b>	MATCH (XP, $\varphi$ )	ALIGNR (XP, $\varphi$ )
[A[B[CD]]]	(A(B(CD)))	((AB)(CD))	W			L	W	
[A[B[CD]]]	(A(B(CD)))	((A(B))(CD))		W		L	W	
[A[[BC]D]]	(A((BC)D))	(A(BC)D)			W	L	W	
						W	L	L

**BINMAX( $\varphi,\omega$ )** prefers winner.

# Support for L.12: *L-branching* → *rebracketed*

Input	Winner	Loser	ALIGNL (XP, $\varphi$ )	BINMIN ( $\varphi, \omega$ )	BINMAX ( $\varphi, br$ )	<b>BINMAX (<math>\varphi, \omega</math>)</b>	MATCH (XP, $\varphi$ )	ALIGNR (XP, $\varphi$ )
[A[B[CD]]]	(A(B(CD)))	((AB)(CD))	W			L	W	
[A[B[CD]]]	(A(B(CD)))	((A(B))(CD))		W		L	W	
[A[[BC]D]]	(A((BC)D))	(A(BC)D)			W	L	W	
						W	L	L

**BINMAX**( $\varphi, \omega$ ) prefers winner.

# Support for L.12: *L-branching* → *rebracketed*

Input	Winner	Loser	ALIGNL (XP, $\varphi$ )	BINMIN ( $\varphi, \omega$ )	BINMAX ( $\varphi, br$ )	<b>BINMAX (<math>\varphi, \omega</math>)</b>	MATCH (XP, $\varphi$ )	ALIGNR (XP, $\varphi$ )
[A[B[CD]]]	(A(B(CD)))	((AB)(CD))	W			L	W	
[A[B[CD]]]	(A(B(CD)))	((A(B))(CD))		W		L	W	
[A[[BC]D]]	(A((BC)D))	(A(BC)D)			W	L	W	
						W	L	L

**BINMAX( $\varphi, \omega$ )** prefers winner.

# Support for L.12: *L-branching* → *rebracketed*

Input	Winner	Loser	ALIGNL (XP,φ)	BINMIN (φ,ω)	BINMAX (φ,br)	<b>BINMAX (φ,ω)</b>	MATCH (XP,φ)	ALIGNR (XP,φ)
[A[B[CD]]]	(A(B(CD)))	((AB)(CD))	W			L	W	
[A[B[CD]]]	(A(B(CD)))	((A(B))(CD))		W		L	W	
[A[[BC]D]]	(A((BC)D))	(A(BC)D)			W	L	W	
						W	L	L

**BINMAX(φ,ω)** prefers winner.

# Support for L.12: *L-branching* → *rebracketed*

Input	Winner	Loser	ALIGNL (XP,φ)	BINMIN (φ,ω)	BINMAX (φ,br)	BINMAX (φ,ω)	<b>MATCH (XP,φ)</b>	ALIGNR (XP,φ)
[A[B[CD]]]	(A(B(CD)))	((AB)(CD))	W			L	W	
[A[B[CD]]]	(A(B(CD)))	((A(B))(CD))		W		L	W	
[A[[BC]D]]	(A((BC)D))	(A(BC)D)			W	L	W	
						W	L	L

**MATCH**(XP,φ) prefers loser;  $[_{XP} \text{ABC}]$  unmatched in winner.

# Support for L.12: *L-branching* → *rebracketed*

Input	Winner	Loser	ALIGNL (XP,φ)	BINMIN (φ,ω)	BINMAX (φ,br)	BINMAX (φ,ω)	<b>MATCH (XP,φ)</b>	ALIGNR (XP,φ)
[A[B[CD]]]	(A(B(CD)))	((AB)(CD))	W			L	W	
[A[B[CD]]]	(A(B(CD)))	((A(B))(CD))		W		L	W	
[A[[BC]D]]	(A((BC)D))	(A(BC)D)			W	L	W	
						W	L	L

**MATCH(XP,φ)** prefers loser;  $[_{XP} ABC]$  unmatched in winner.

# Support for L.12: *L-branching* → *rebracketed*

Input	Winner	Loser	ALIGNL (XP,φ)	BINMIN (φ,ω)	BINMAX (φ,br)	BINMAX (φ,ω)	MATCH (XP,φ)	<b>ALIGNR (XP,φ)</b>
[A[B[CD]]]	(A(B(CD)))	((AB)(CD))	W			L	W	
[A[B[CD]]]	(A(B(CD)))	((A(B))(CD))		W		L	W	
[A[[BC]D]]	(A((BC)D))	(A(BC)D)			W	L	W	
						W	L	L

ALIGNR(XP,φ) prefers loser; **C]→C** in winner.



# Support for L.12: *mixed-branching* → *isomorphic*

Input	Winner	Loser	ALIGNL (XP, $\varphi$ )	BINMIN ( $\varphi, \omega$ )	<b>BINMAX</b> ( $\varphi, br$ )	BINMAX ( $\varphi, \omega$ )	MATCH (XP, $\varphi$ )	ALIGNR (XP, $\varphi$ )
[A[B[CD]]]	(A(B(CD)))	((AB)(CD))	W			L	W	
[A[B[CD]]]	(A(B(CD)))	((A(B))(CD))		W		L	W	
					W	L	W	
[[[AB]C]D]	((AB)(CD))	((([AB]C)D))				W	L	L

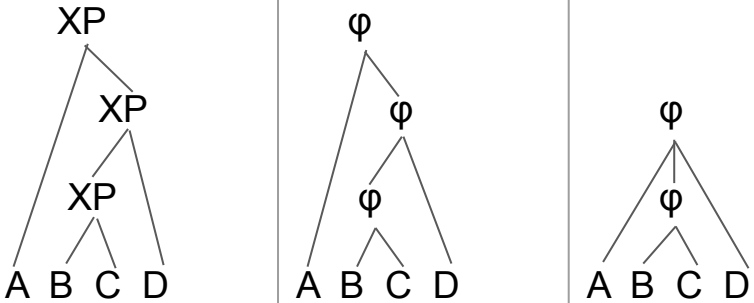
**BINMAX( $\varphi, branches$ )** prefers winner; loser contains **ternary** ( $\varphi A \varphi D$ ).

# Support for L.12: *mixed-branching* → *isomorphic*

Input	Winner	Loser	ALIGNL (XP, $\varphi$ )	BINMIN ( $\varphi, \omega$ )	BINMAX ( $\varphi, br$ )	<b>BINMAX (<math>\varphi, \omega</math>)</b>	MATCH (XP, $\varphi$ )	ALIGNR (XP, $\varphi$ )
[A[B[CD]]]	(A(B(CD)))	((AB)(CD))	W			L	W	
[A[B[CD]]]	(A(B(CD)))	((A(B))(CD))		W		L	W	
					W	L	W	
[[[AB]C]D]	((AB)(CD))	((([AB]C)D))				W	L	L

**BINMAX( $\varphi, \omega$ )** prefers loser.

# Support for L.12: *mixed-branching* → *isomorphic*

Input	Winner	Loser	ALIGNL (XP,φ)	BINMIN (φ,ω)	<b>BINMAX (φ,br)</b>	<b>BINMAX (φ,ω)</b>	<b>MATCH (XP,φ)</b>	ALIGNR (XP,φ)
[A[B[CD]]]	(A(B(CD)))	((AB)(CD))	W			L	W	
[A[B[CD]]]	(A(B(CD)))	((A(B))(CD))		W		L	W	
					W	L	W	
[[[AB]C]D]	((AB)(CD))	((((AB)C)D))				W	L	L

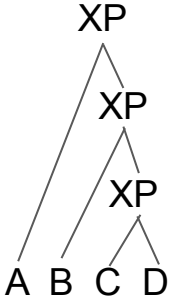
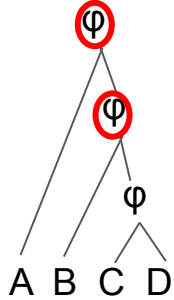
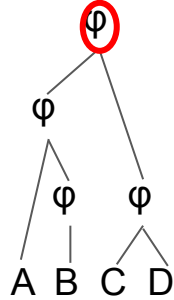
**MATCH(XP,φ)** prefers winner, but we already know **BINMAX(φ,ω) >> MATCH(XP,φ)**

# Support for L.12: *R-branching* → *isomorphic*

Input	Winner	Loser	ALIGNL (XP,φ)	<b>BINMIN</b> (φ,ω)	BINMAX (φ,br)	BINMAX (φ,ω)	MATCH (XP,φ)	ALIGNR (XP,φ)
[A[B[CD]]]	(A(B(CD)))	((AB)(CD))	W			L	W	
				W		L	W	
[A[[BC]D]]	(A((BC)D))	(A(BC)D)			W	L	W	
[[[AB]C]D]	((AB)(CD))	((((AB)C)D))				W	L	L

**BINMIN**(φ,ω) prefers winner; loser contains **unary** ( $\varphi$ B).

# Support for L.12: *R-branching* $\rightarrow$ *isomorphic*

Input	Winner	Loser	ALIGNL (XP, $\varphi$ )	BINMIN ( $\varphi,\omega$ )	BINMAX ( $\varphi,br$ )	<b>BINMAX (<math>\varphi,\omega</math>)</b>	MATCH (XP, $\varphi$ )	ALIGNR (XP, $\varphi$ )
[A[B[CD]]]	(A(B(CD)))	((AB)(CD))	W			L	W	
				W		L	W	
[A[[BC]D]]	(A((BC)D))	(A(BC)D)			W	L	W	
[[[AB]C]D]	((AB)(CD))	((((AB)C)D))				W	L	L

**BINMAX( $\varphi,\omega$ )** prefers loser.

# Support for L.12: *R-branching* $\rightarrow$ *isomorphic*

Input	Winner	Loser	ALIGNL (XP, $\varphi$ )	BINMIN ( $\varphi,\omega$ )	BINMAX ( $\varphi,br$ )	BINMAX ( $\varphi,\omega$ )	MATCH (XP, $\varphi$ )	ALIGNR (XP, $\varphi$ )
<p>A parse tree for the input [A[B[CD]]]. The root node is XP (blue). It has three children: A (blue), B (red), and another XP (red). This second XP node has two children: C (blue) and D (blue). The tree structure is XP(A, XP(B, C, D)).</p>	<p>A parse tree for the input (A(B(CD))). The root node is <math>\varphi</math>. It has three children: A, B, and another <math>\varphi</math>. This second <math>\varphi</math> node has two children: C and D. The tree structure is <math>\varphi(A, \varphi(B, C, D))</math>.</p>	<p>A parse tree for the input ((A(B))(CD)). The root node is <math>\varphi</math>. It has two children: another <math>\varphi</math> and another <math>\varphi</math>. The first <math>\varphi</math> node has two children: A and B. The second <math>\varphi</math> node has two children: C and D. The tree structure is <math>\varphi(\varphi(A, B), \varphi(C, D))</math>.</p>	W			L	W	
[A[B[CD]]]	(A(B(CD)))	((A(B))(CD))		W		L	W	
[A[[BC]D]]	(A((BC)D))	(A(BC)D)			W	L	W	
[[[AB]C]D]	((AB)(CD))	((((AB)C)D))				W	L	L

ALIGNL(XP, $\varphi$ ) prefers winner; **[B $\rightarrow$ (B** in loser.

# Support for L.12: *R-branching* $\rightarrow$ *isomorphic*

Input	Winner	Loser	ALIGNL (XP, $\varphi$ )	BINMIN ( $\varphi,\omega$ )	BINMAX ( $\varphi,br$ )	<b>BINMAX (<math>\varphi,\omega</math>)</b>	MATCH (XP, $\varphi$ )	ALIGNR (XP, $\varphi$ )
			W			L	W	
[A[B[CD]]]	(A(B(CD)))	((A(B))(CD))		W		L	W	
[A[[BC]D]]	(A((BC)D))	(A(BC)D)			W	L	W	
[[[AB]C]D]	((AB)(CD))	((((AB)C)D))				W	L	L

**BINMAX( $\varphi,\omega$ )** prefers loser.

**S.Msp.Mps**

Pure MATCH



Can we get the pattern without ALIGN?

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- No!

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(2) CON.Msp.Mps

(a) Mapping constraints:

(i) MATCH(XP, $\varphi$ )

(ii) MATCH( $\varphi$ ,XP): Assign one violation for every node of category  $\varphi$  in the prosodic tree such that there is no node of category XP in the syntactic tree that dominates all and only the same terminal nodes as  $\varphi$ .

(b) Markedness constraints:

(i) BINMIN( $\varphi$ , $\omega$ )

(ii) BINMAX( $\varphi$ , $\omega$ )

(iii) BINMAX( $\varphi$ ,branches)

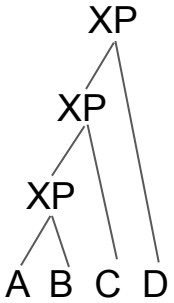
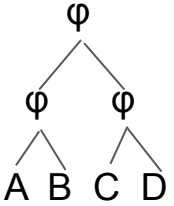
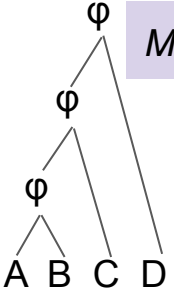
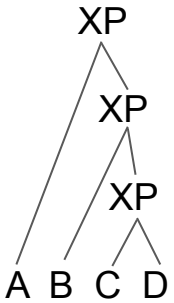
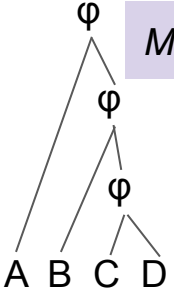
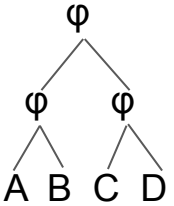
# The Asymmetry Problem

Input	Winner	Loser	MATCH (XP, $\varphi$ )	MATCH ( $\varphi$ ,XP)	BINMAX ( $\varphi$ , $\omega$ )	BINMIN ( $\varphi$ , $\omega$ )	BINMAX ( $\varphi$ ,br)
			L	L	W	e	e
			W	W	L	e	e

# The Asymmetry Problem

Input	Winner	Loser	MATCH (XP, $\varphi$ )	MATCH ( $\varphi$ ,XP)	<b>BINMAX</b> ( $\varphi,\omega$ )	BINMIN ( $\varphi,\omega$ )	BINMAX ( $\varphi,br$ )
			L	L	W	e	e
			W	W	L	e	e

# The Asymmetry Problem

Input	Winner	Loser	MATCH (XP, $\varphi$ )	MATCH ( $\varphi$ ,XP)	BINMAX ( $\varphi$ , $\omega$ )	BINMIN ( $\varphi$ , $\omega$ )	BINMAX ( $\varphi$ ,br)
	<p><i>Mismatch</i></p> 	<p><i>Match</i></p> 	L	L	W	e	e
	<p><i>Match</i></p> 	<p><i>Mismatch</i></p> 	W	W	L	e	e



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  - Hypothetical STRONGEND doesn't work either (see (26) in our paper).
- *Conclusion*: Unless we find a plausible asymmetric markedness constraint, **we need ALIGN constraints.\***

\*Or some other, as yet undiscovered, asymmetric mapping constraint.

**S.Asp.Aps**

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ALIGNL( $\varphi$ ,XP)

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ALIGNR(XP, $\varphi$ )                  ALIGNR( $\varphi$ ,XP)

(b) Markedness constraints

same binarity constraints as before

# The Ambivalence Problem

Input	Winner	Loser	ALIGNL (XP, $\varphi$ )	ALIGNR (XP, $\varphi$ )	ALIGNL ( $\varphi$ ,XP)	ALIGNR ( $\varphi$ ,XP)	BINMAX ( $\varphi$ , $\omega$ )	BINMIN ( $\varphi$ , $\omega$ )	BINMAX ( $\varphi$ ,br)
			e	e	e	e	e	e	e
			e	e	e	e	e	e	e

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- **No markedness constraint**, standard or novel, could solve this problem; difference between  $[A[[BC]D]] \rightarrow (A((BC)D))$  and  $[[A[BC]]D] \rightarrow ((A(BC))D)$  comes down to **mapping (faithfulness)**.

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- The mapping constraint WRAP(XP) can't solve the problem; all of the outputs in our systems satisfy it perfectly.
- So we need a **MATCH** constraint.

# Conclusion

Using SPOT, we have found that Japanese  $\phi$ -phrasing involves **Match and Align**.

Direction for future research:

- Expanding past 4-word phrases, our systems make predictions for larger prosodic structures. Are these borne out?
- Are the other languages in the factorial typology of S.Msp.Asp empirically supported? (e.g. mirror-image Japanese?)